



# A biased random-key genetic algorithm for a prize-collecting directed Steiner forest network design problem



## Summary

- Prize collecting directed k-hop Steiner forest (PCk-HSF) problem
- Wireless backhaul network planning as a PCk-HSF problem with additional constraints
- Biased random-key genetic algorithm (BRKGA) for wireless backhaul network planning focusing on the decoder
- Application of the BRKGA to a "real" instance of wireless backhaul network planning
- Concluding remarks

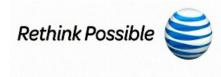


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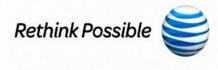
- Let G = (V,E) be a given directed graph
- Let  $V_r$ ,  $V_s$ , and  $V_d$  partition the node set, i.e.

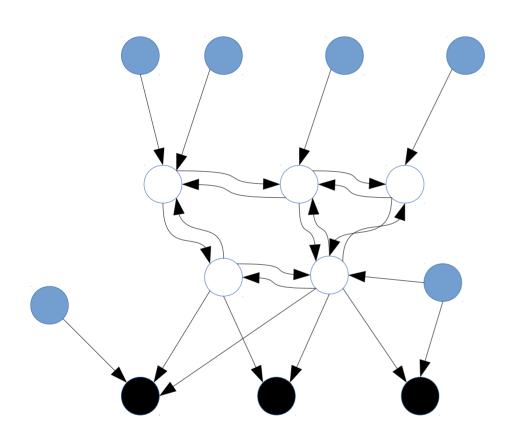
$$V_r \cup V_s \cup V_d = V$$
 and  $V_r \cap V_s \cap V_d = \emptyset$ 

- $V_r$  set of root nodes
- V<sub>s</sub> set of Steiner nodes
- $V_d$  set of demand nodes



- Let G = (V,E) be a given directed graph
- Arcs are directed
  - From demand nodes to Steiner nodes and root nodes
  - From Steiner nodes to root nodes
- For each pair of nodes v, u ∈ V<sub>s</sub>
  - $-(u,v) \in E$
  - $-(v,u) \in E$





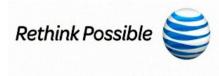


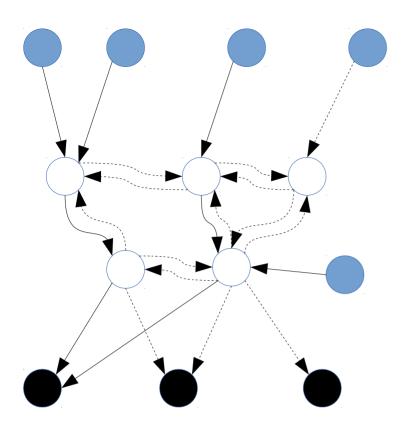
Steiner node

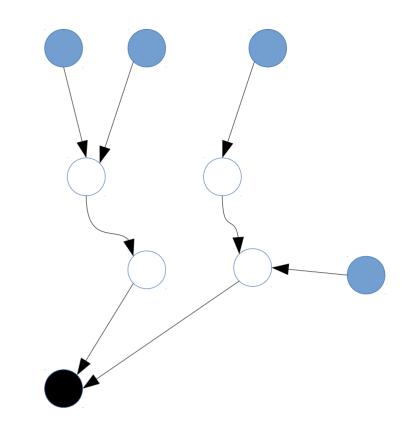
Root node

- A Directed Steiner Tree  $T = (V[T] \subseteq V, E[T] \subseteq E)$  with root in  $r \in V_r[T]$  is a loopless connected subgraph of G with a unique root r.
  - For each demand node u ∈ V[T] there is a unique path in T
     from u to r.

- A k-Hop Directed Steiner Tree is a directed Steiner tree
   (connected subgraph of G with a unique root r) such that:
  - Any path from a demand node u to the root r has no more than k+2 nodes, including u and r.
- A k-Hop Directed Steiner Forest is a collection of disjoint k-Hop Directed Steiner Trees



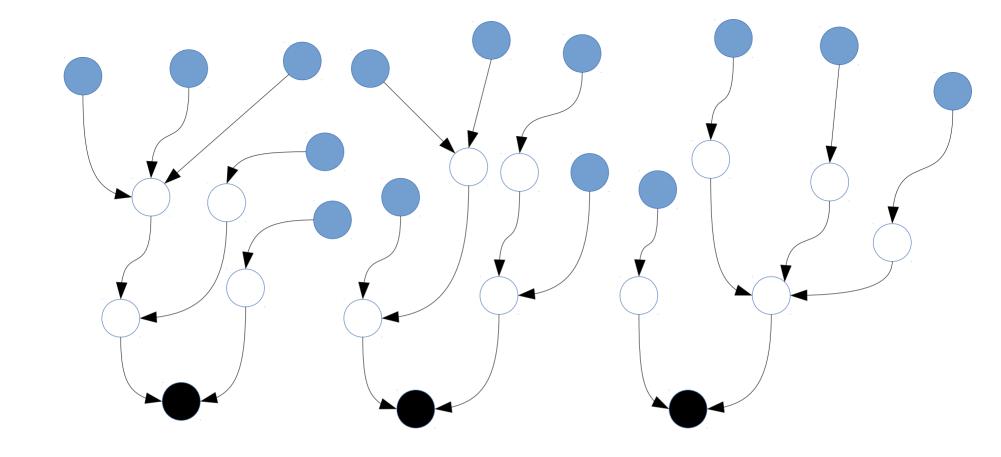




Acyclic directed graph

2-Hop Directed Steiner tree

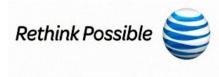




#### 2-Hop Directed Steiner Forest



- A demand function d: V<sub>d</sub> → R<sup>+</sup> gives the prize to be collected from each demand node if it is a leaf on the tree.
- A cost function c:  $V_s \rightarrow \mathbb{R}^+$  gives the cost of each Steiner vertex.



- Given a directed graph G = (V,E), vertex partition  $V_r$ ,  $V_s$ , and  $V_d$ , hop parameter k, demand function d, and cost function c
- FIND: A k-hop directed Steiner forest F such that the profit

sum [ $v \in V_d[F]$ ]  $d_v$  — sum [ $v \in V_s[F]$ ]  $c_v$  is maximized.

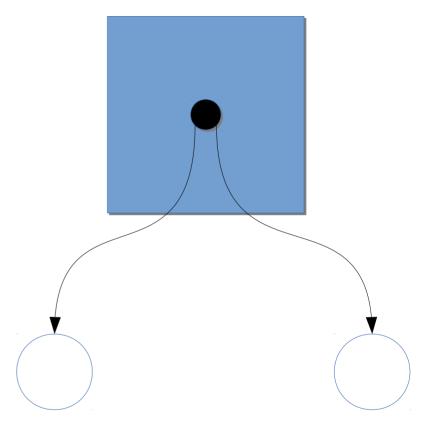


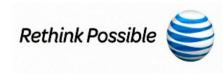
Given a geographical region where locations are represented as lat-long coordinates, where

- V<sub>d</sub> is the set of nodes representing traffic origination points (demand points) in the region
- V<sub>s</sub> is the set of nodes representing locations where equipment for traffic collection and routing is located (e.g. utility poles)
- $-V_r$  is the set of nodes representing fibered access points (FAP), e.g. remote terminals (RT), macrocell (MC), or central offices.

#### Constraints: Demand splitting

- Estimate for each block total demand is placed in center of block
- Block can be served by one or more antennae so demand can be split among them
- Because of this undirected cycles can be introduced resulting in a directed acyclic graph (DAG)





#### Constraints: Access equipment action radii & capacities

#### Action radii

- Wi-Fi: 100 m

- LTE: 400 m

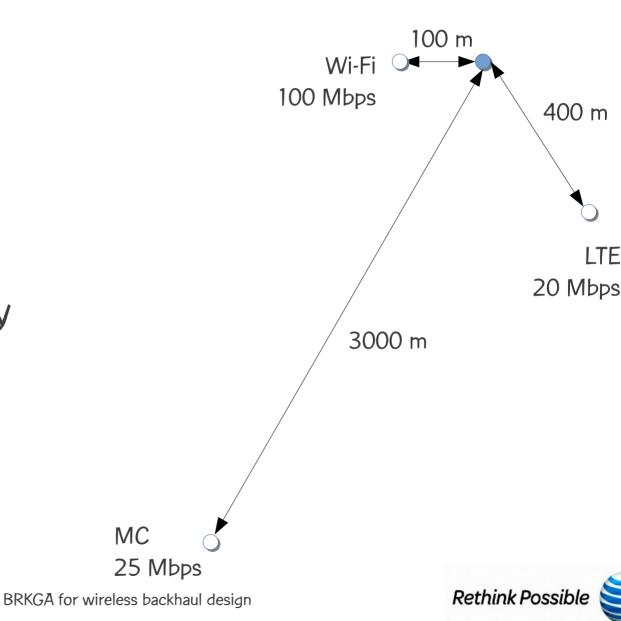
macrocell: 3 km

#### Processing capacity

Wi-Fi: 100 Mbps

LTE: 20 Mbps

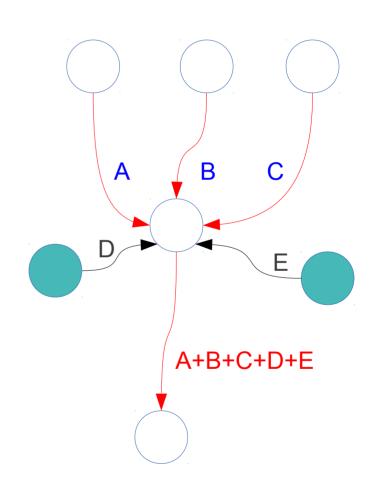
macrocell: 25 Mbps



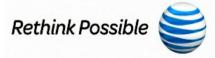
ITF

#### Constraints: Retransmitter equipment capacity

- Action radii: 1 km
- Processing capacity
  - Flow that equipment receives from other wireless retransmitters plus flow it sends to other retransmitters is limited to at most 100 Mbps



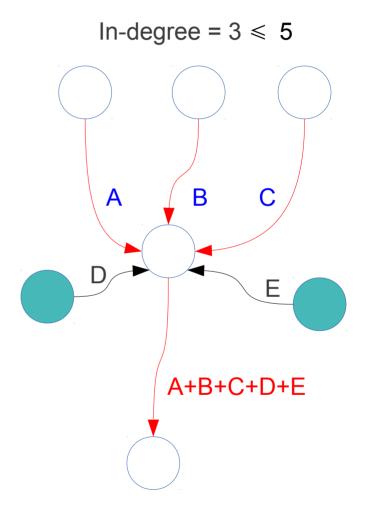
A+B+C+A+B+C+D+E ≤100



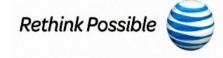
#### Constraints: Retransmitter equipment capacity

#### Fan-in constraint

 A limited number (5) of neighboring transmission equipment can flow traffic into equipment



A+B+C+A+B+C+D+E ≤100

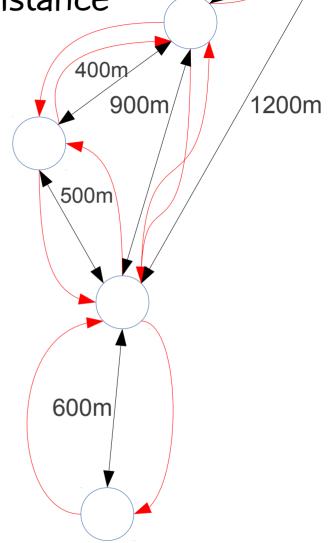


Constraints: Line of sight and minimum distance

PCSF Problem assumed, for each pair  $u, v \in V_s$  that  $(u,v) \in E$  and  $(v,u) \in E$ 

Not so in wireless backhaul planning:  $(u,v) \in E$  and  $(v,u) \in E$  only if

 Poles u and v are within 1000 m of each other



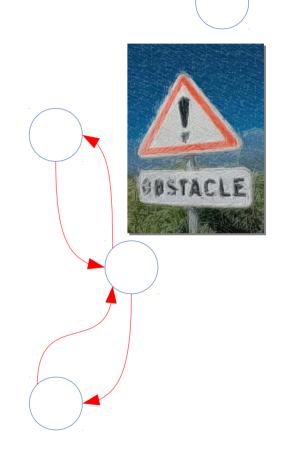
400m

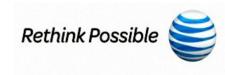
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Not so in wireless backhaul planning:  $(u,v) \in E$  and  $(v,u) \in E$  only if

- Poles u and v are within 1000 m of each other
- Poles u and v are in each other's line of sight

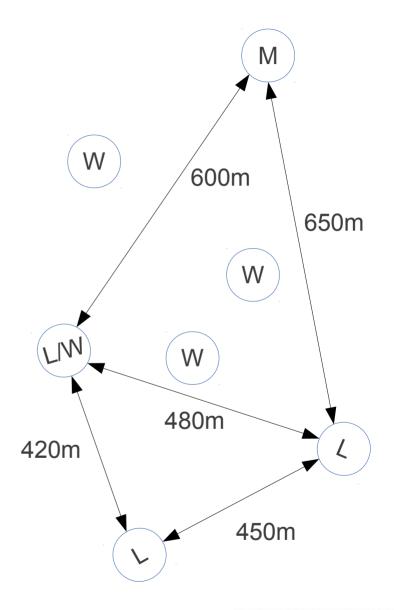


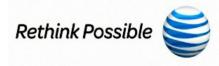


Constraints: Interference

LTE and MC use licensed spectrum and can interfere with each other.

- Pairs of LTE antennae must be separated by at least 400m
- LTE and macrocells by at least 500m
- No constraint on Wi-Fi exists since it uses non-licensed spectrum





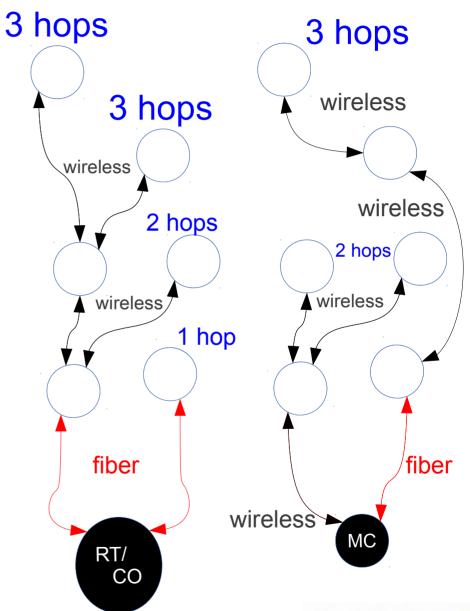
Constraints: k-Hops

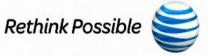
First hop is from FAP (root) to pole

- If root is Central Office (CO) or RT link must use fiber
- If root is macrocell link can be fiber or wireless

All other links are wireless

Number of hops is limited to k = 2 or 3 (in case first link is fiber)

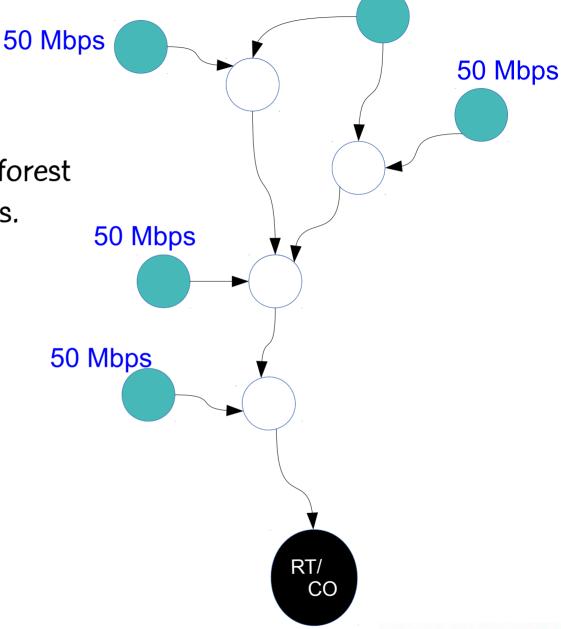


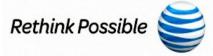


Constraints: Traffic flow

Traffic that is backhauled from demand points to root nodes of forest is limited by equipment capacities.

Only traffic that reaches roots is counted as revenue.



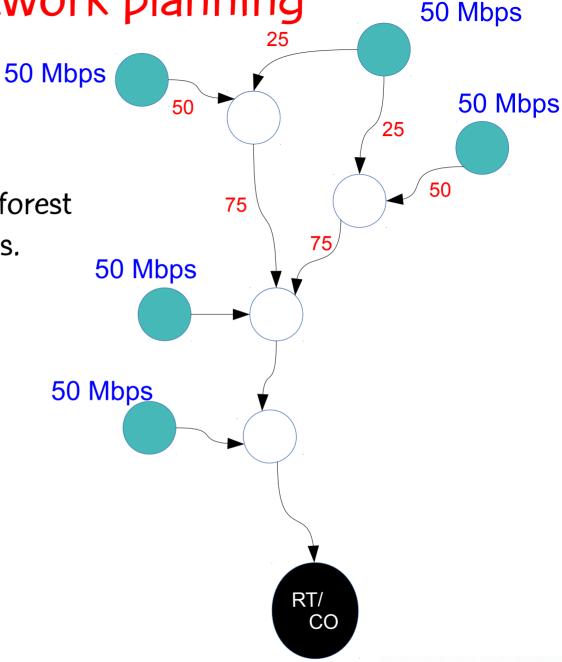


50 Mbps

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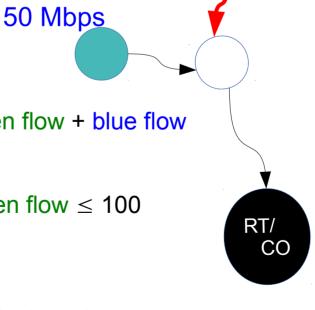
Constraints: Traffic flow

Traffic that is backhauled from demand points to root nodes of forest is limited by equipment capacities.

Only traffic that reaches roots is counted as revenue.

Flow conservation: red flow = cyan flow + green flow + blue flow

Capacity constraint: red flow + cyan flow + green flow ≤ 100

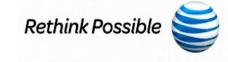


75

75

50

50 Mbps



50 Mbps

25

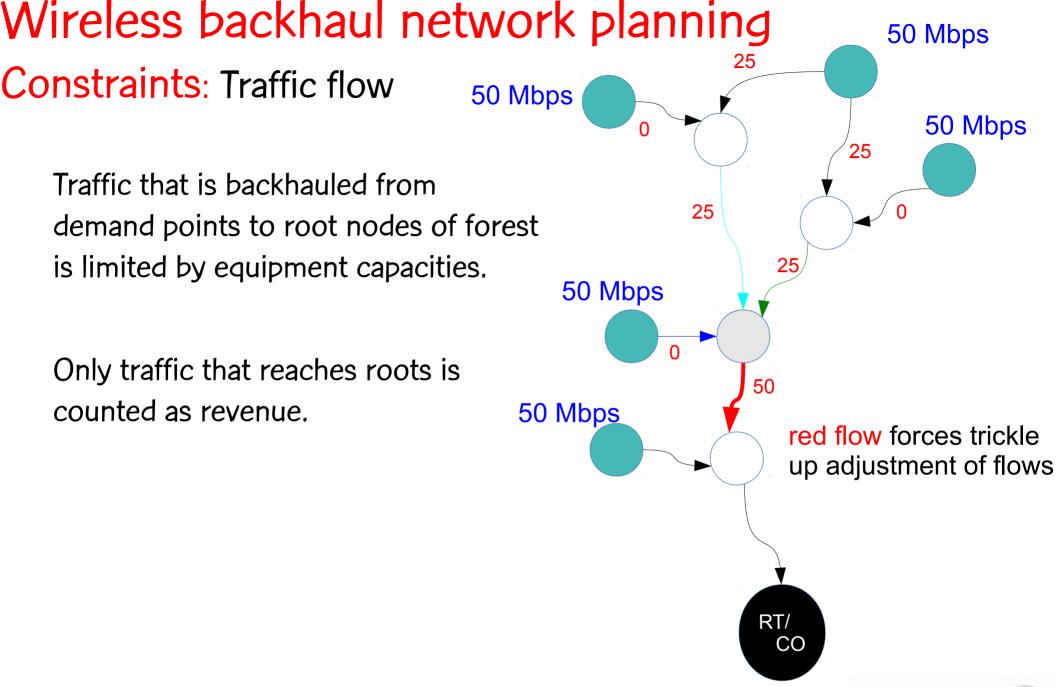
50 Mbps

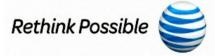
50 Mbps

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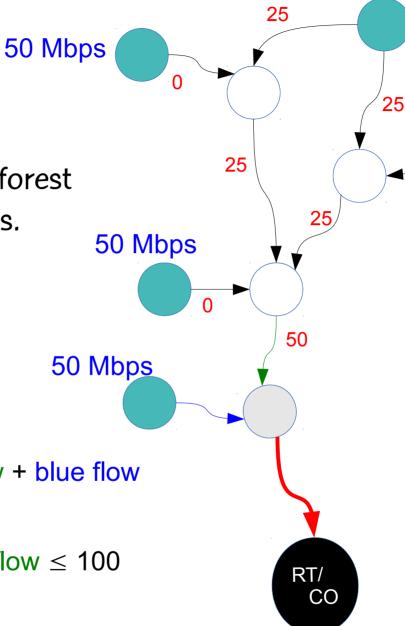
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Flow conservation: red flow = green flow + blue flow

Capacity constraint: red flow + green flow ≤ 100





50 Mbps

50 Mbps

Constraints: Traffic flow

Traffic that is backhauled from demand points to root nodes of forest is limited by equipment capacities.

Only traffic that reaches roots is counted as revenue.

50 Mbps

red flow forces trickle up adjustment of flows

75

RT/ CO

Flow conservation: red flow = green flow + blue flow = 75

Capacity constraint: red flow + green flow = 100 ≤ 100



50 Mbps

12.5

50 Mbps

50 Mbps

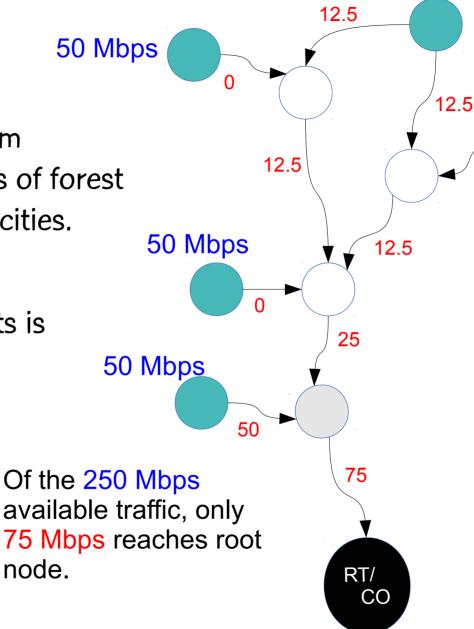
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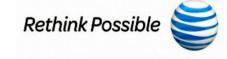
12.5

Wireless backhaul network planning Constraints: Traffic flow

Traffic that is backhauled from demand points to root nodes of forest is limited by equipment capacities.

Only traffic that reaches roots is counted as revenue.





50 Mbps

50 Mbps

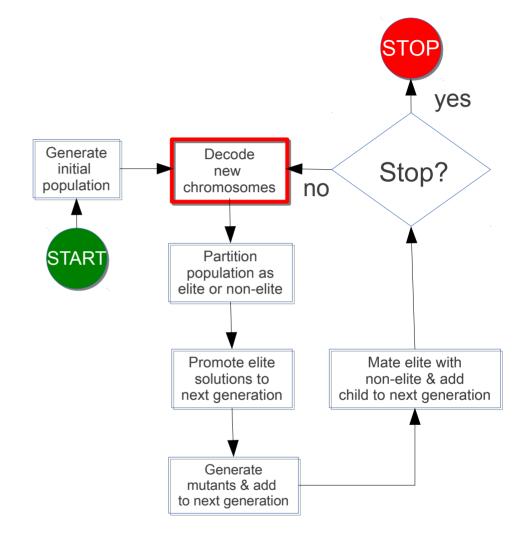
node.

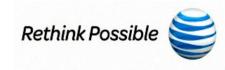
# Biased random-key genetic algorithm

Gonçalves & Resende (J. of Heuristics, 2011)

#### Features of a BRKGA

- Fixed chromosome encoding using a vector of random keys
- Well-defined problem-independent evolutionary process for vectors of keys
  - parametrized uniform crossover
  - mutants in place of mutation
  - elitism
  - bias in selection of mates and in crossover
- Decoder is only problem-dependent component





## Decoder for wireless backhaul planning

- Biased Random-Key Genetic Algorithm
  - Learn the best network layout and equipment placement
- A solution is encoded by a vector x ∈ [0,1]<sup>n</sup>
  - Where  $n = 5 \times \#$  of poles

Activation order

Activation parameters

Evaluation order

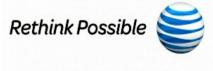
Neighborhood evaluation order

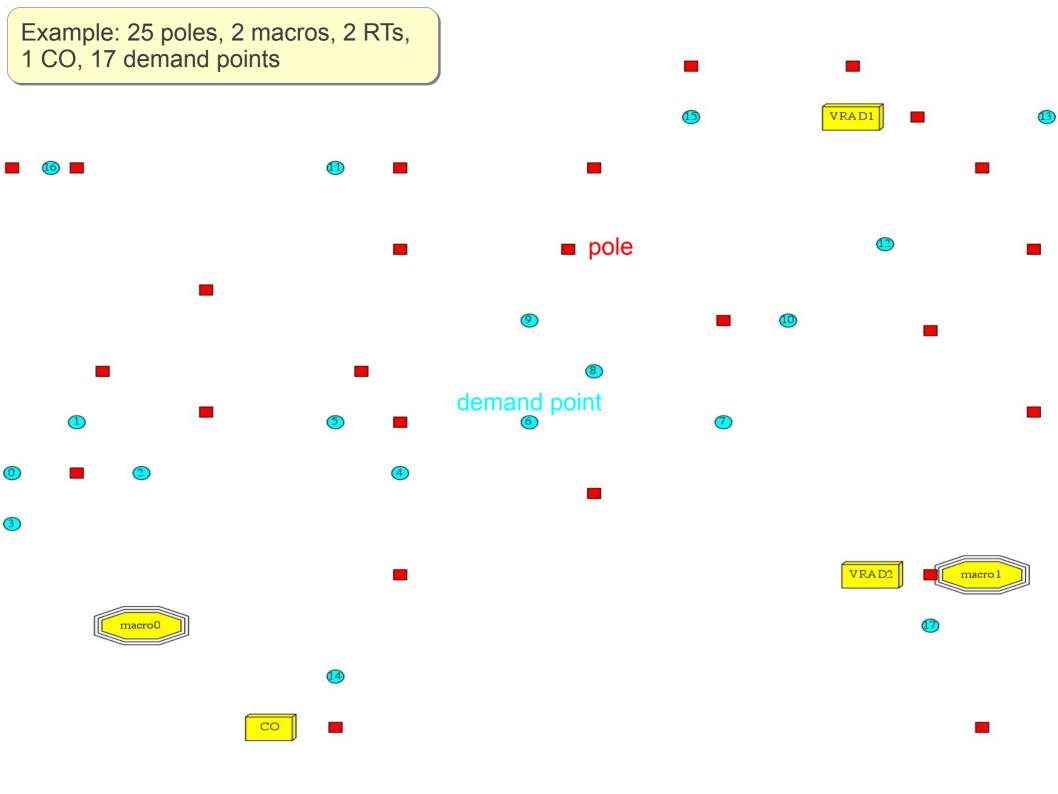
Minimum tree level of pole

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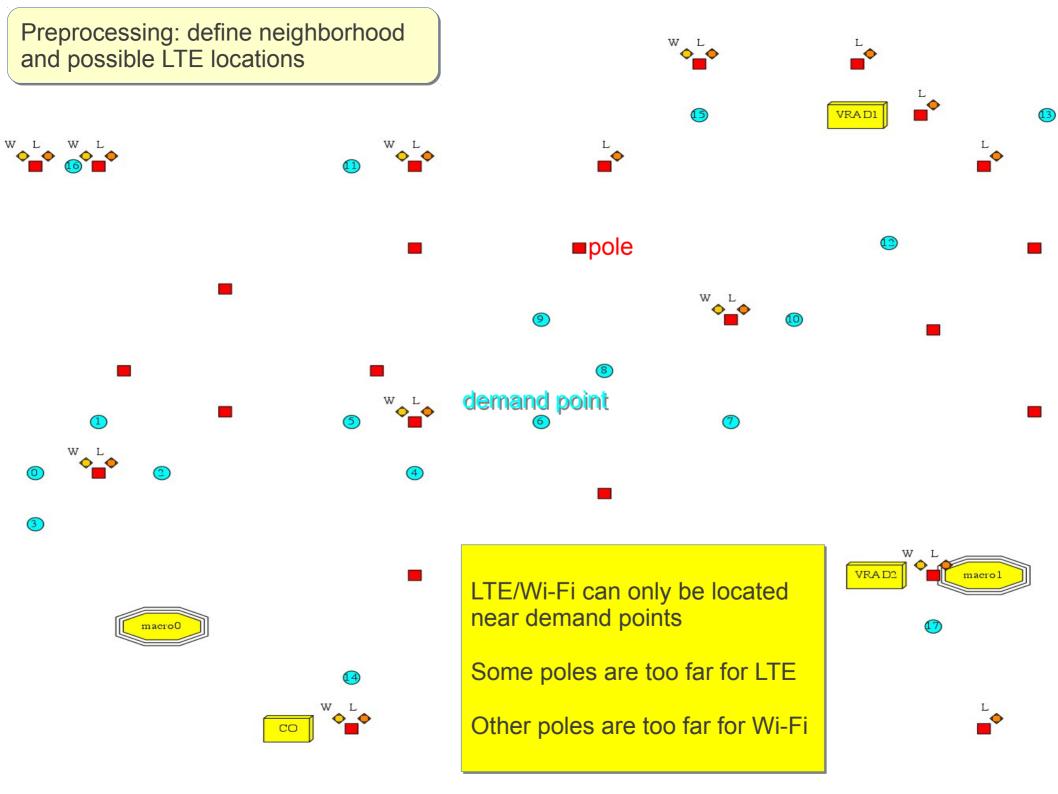
#### Decoder

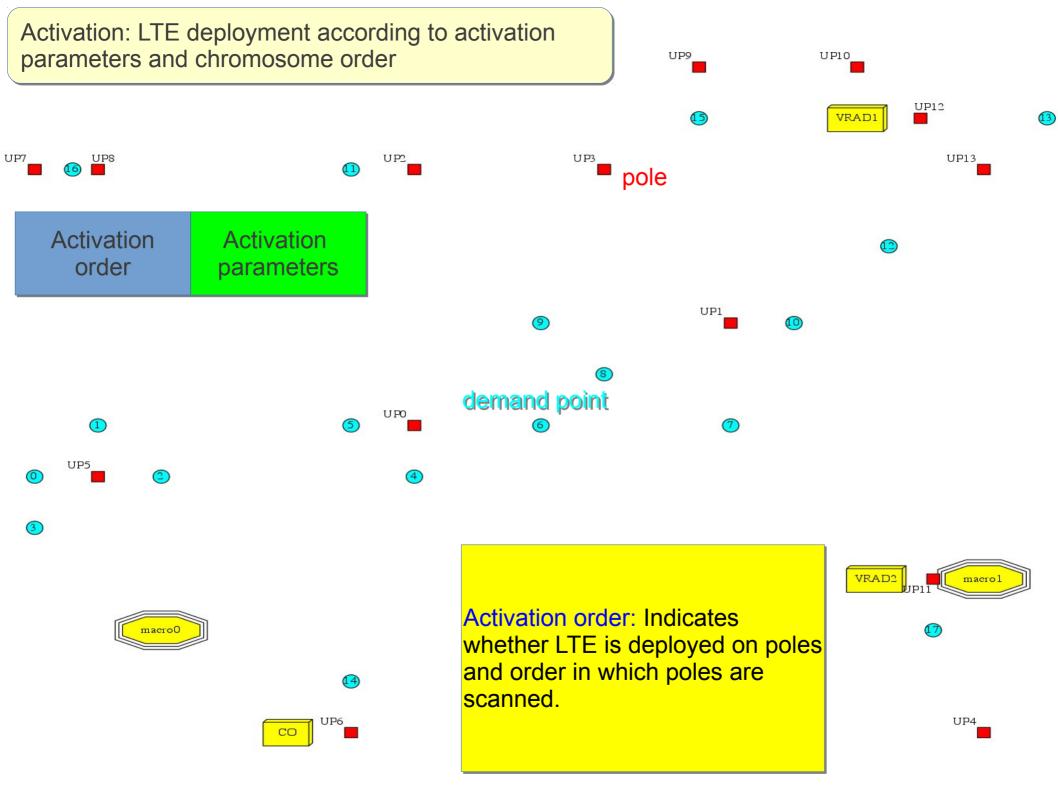
- Define activation order & install LTE on poles
- Build backhaul graph
- Remove unused equipment
- Compute maximum flow from demand points to FAPs
- Remove unused equipment & poles
- Compute cost and revenue and return objective function value

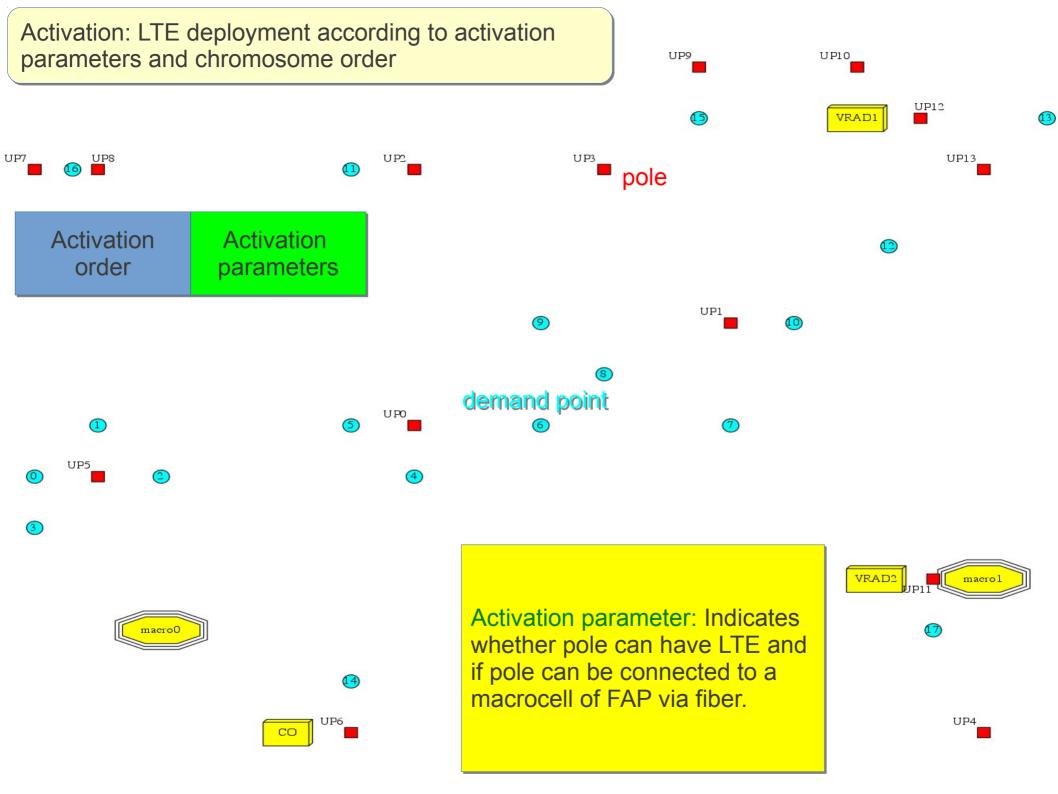


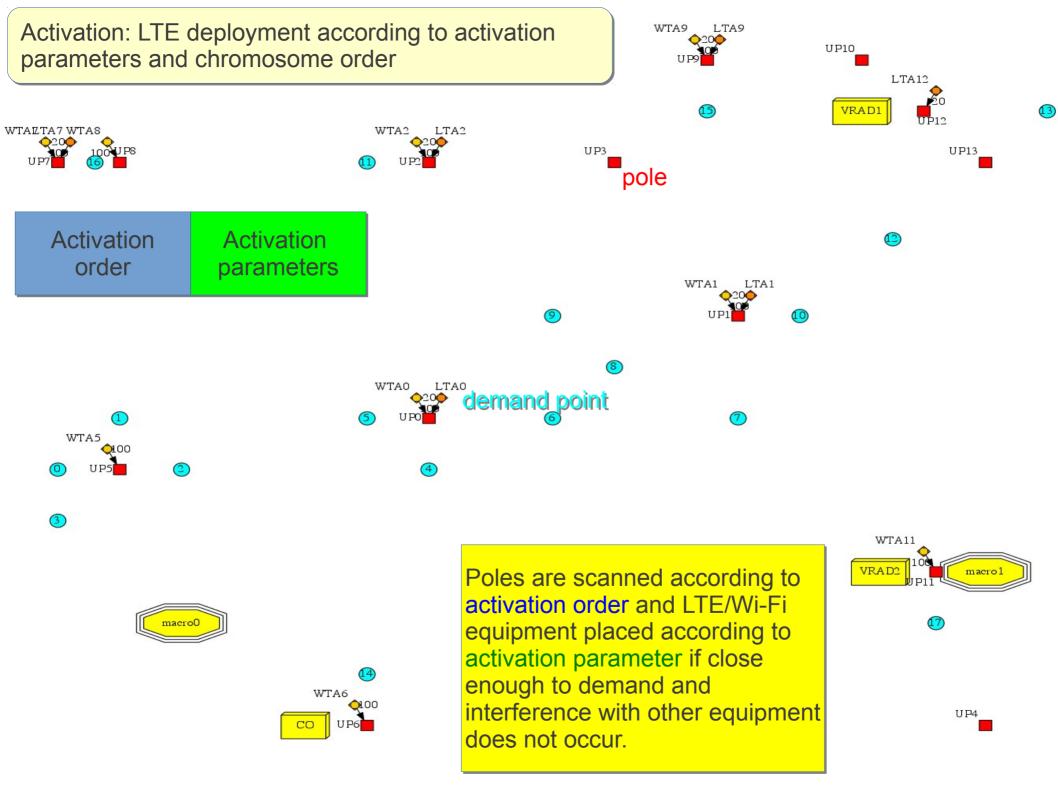


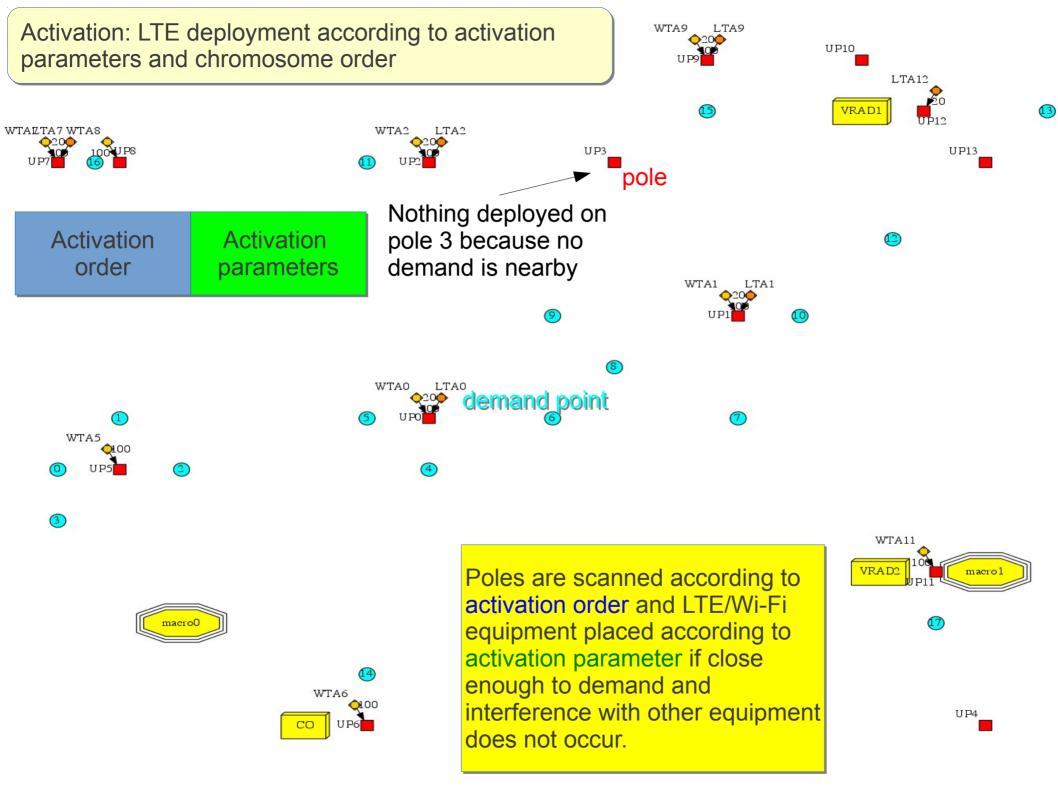
# Preprocessing: define neighborhood and possible LTE locations VRA D1 pole demand point macro1 macro0

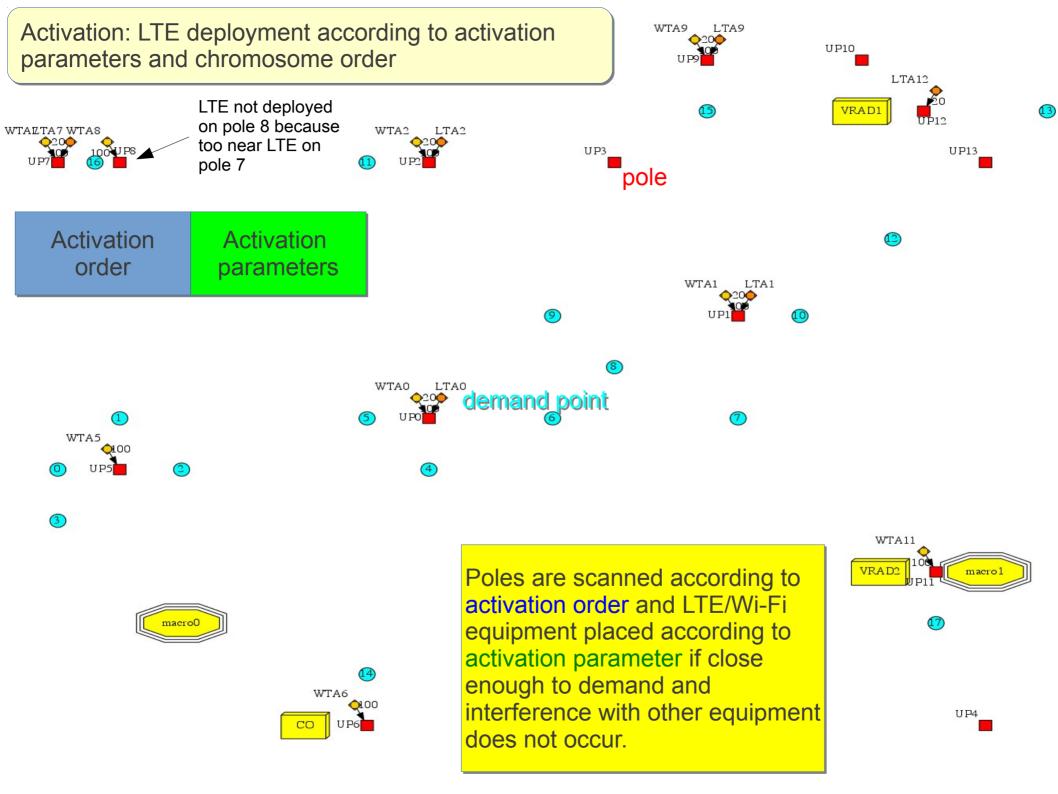


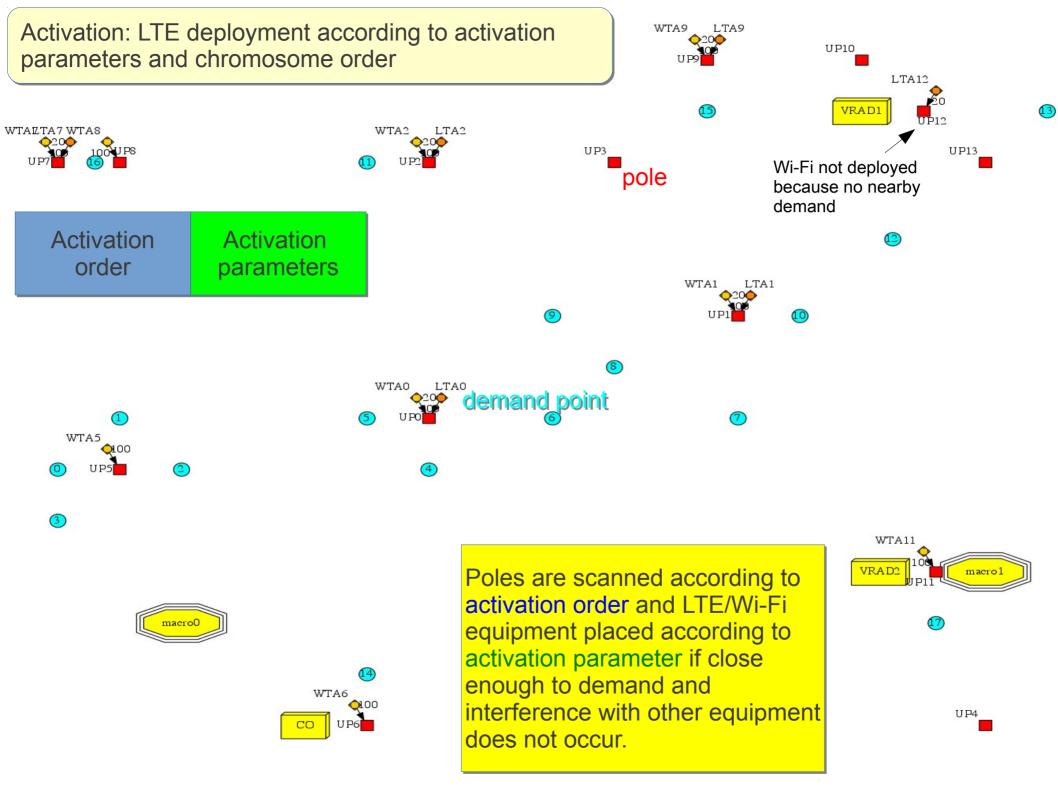




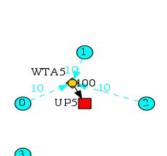




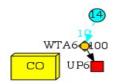


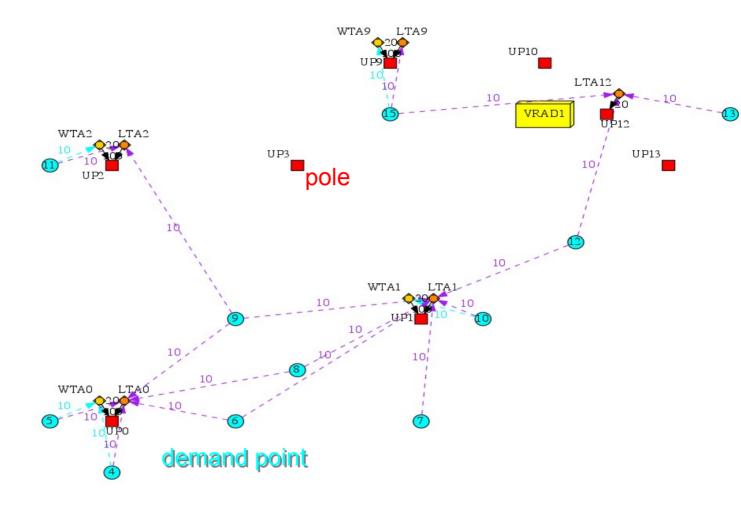


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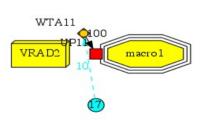




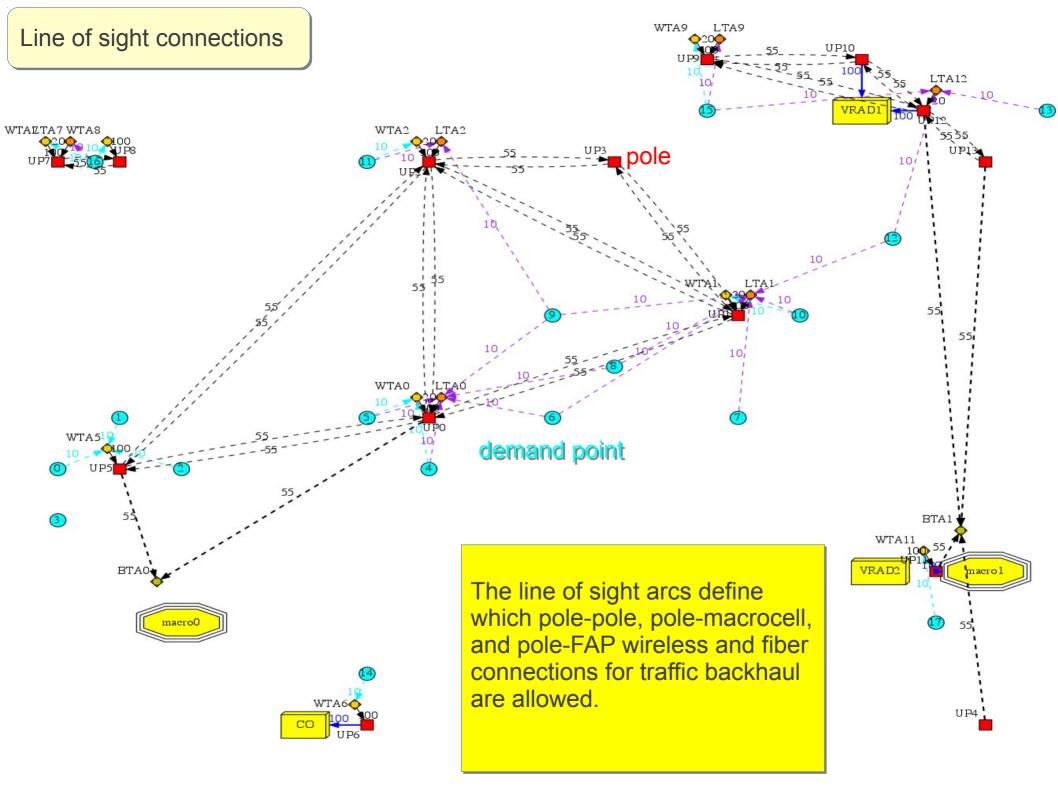


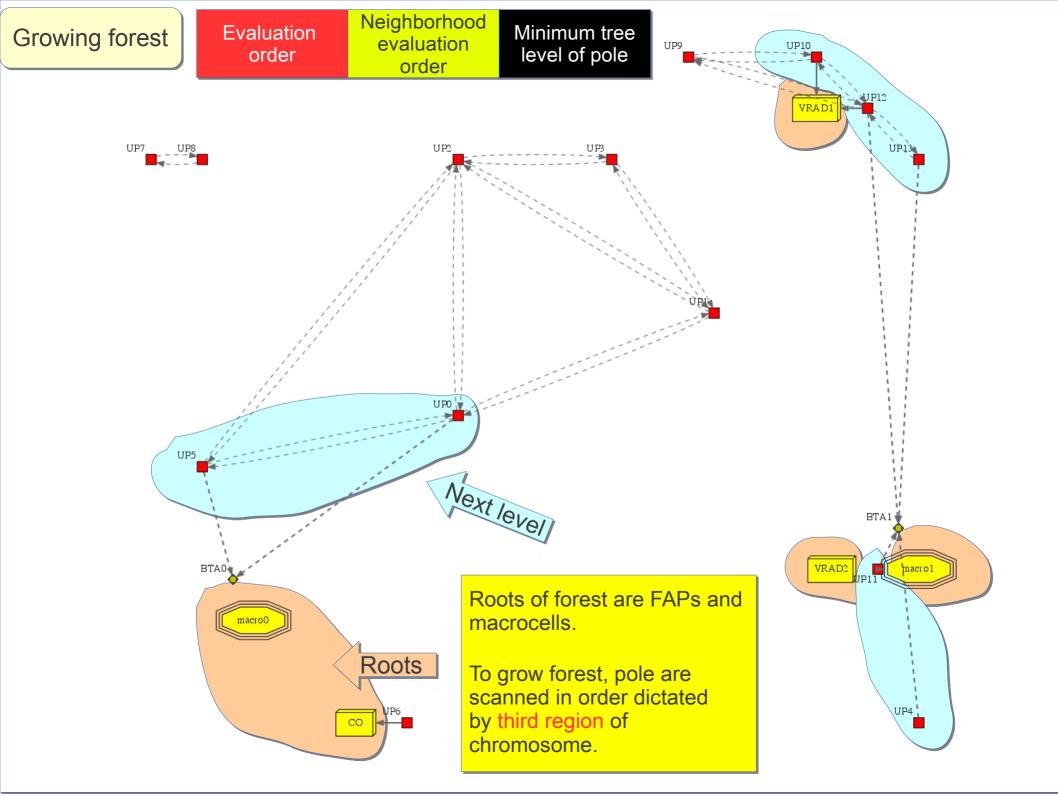
Demands are homed-in on poles that are within range.

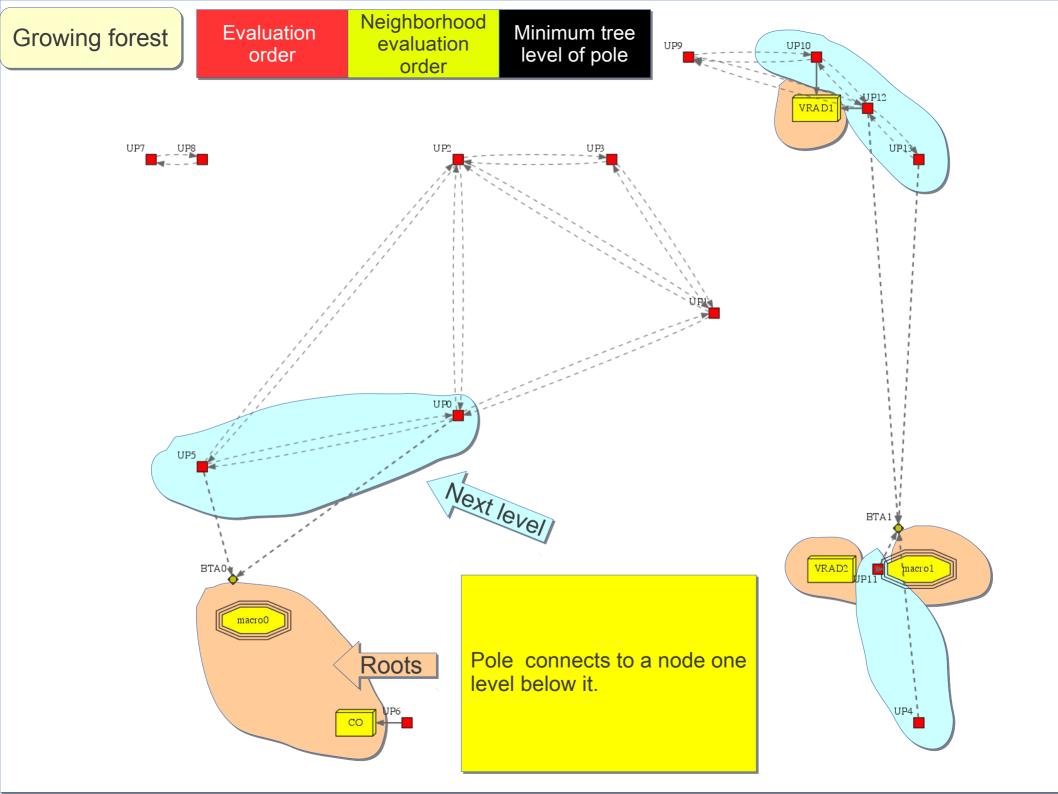
If demand is near more than one pole, its traffic is allowed to be split among these poles.

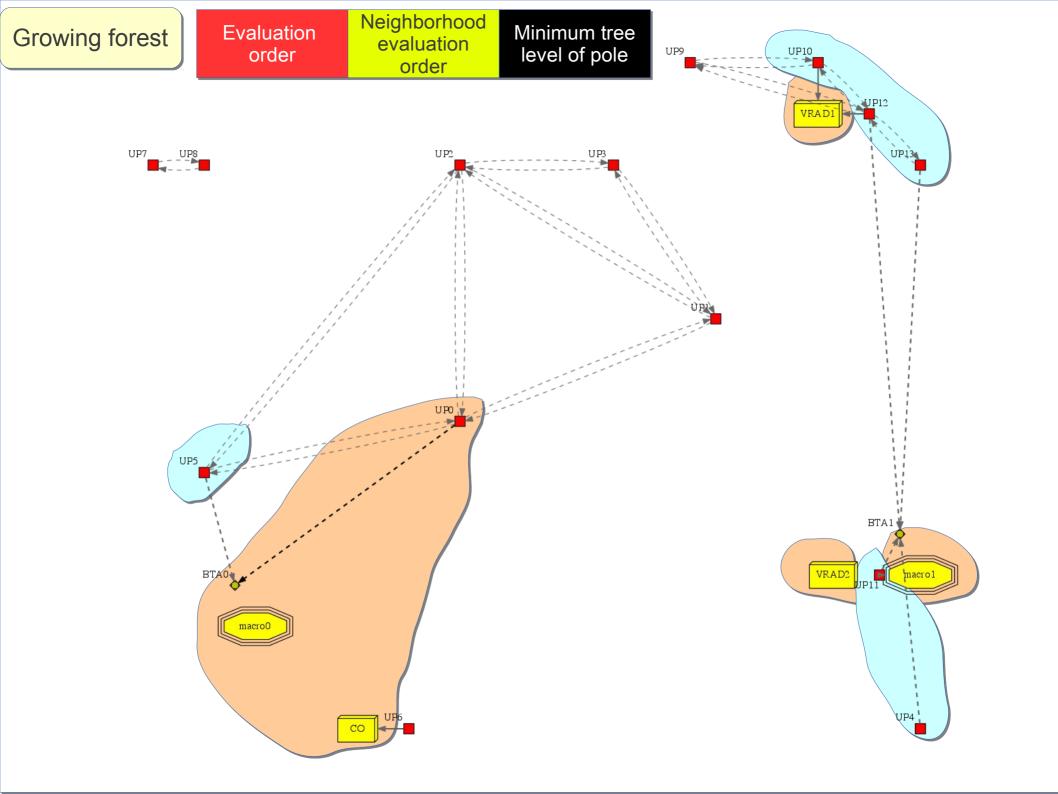


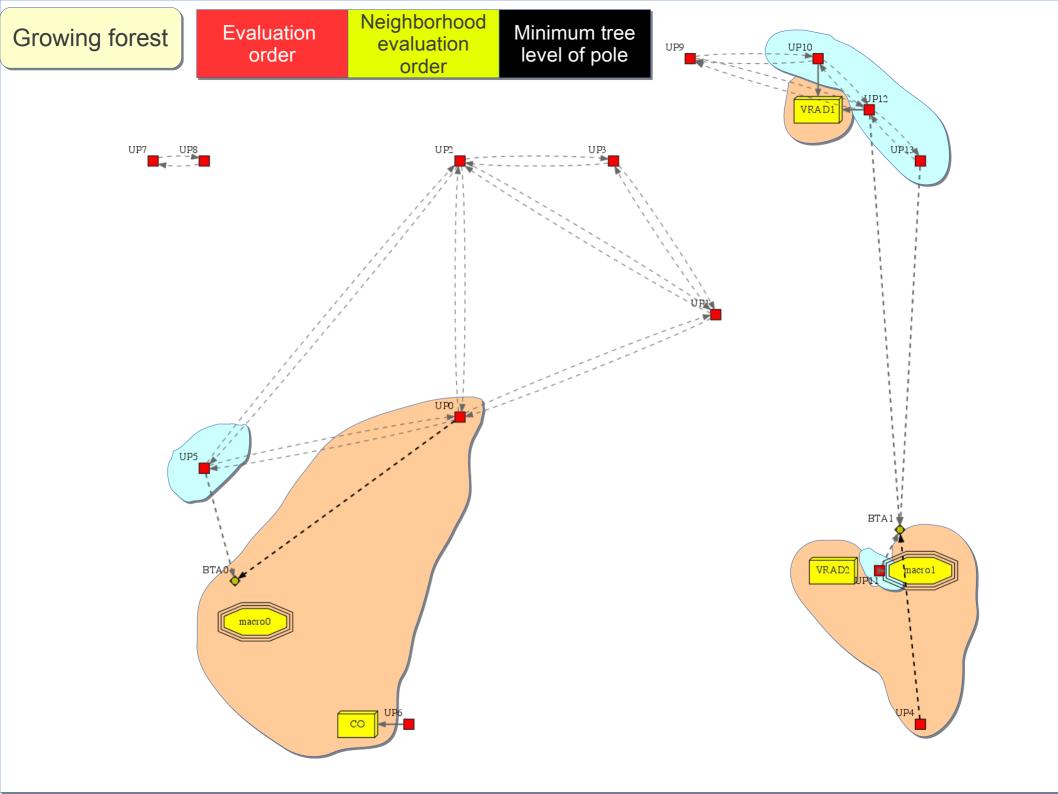


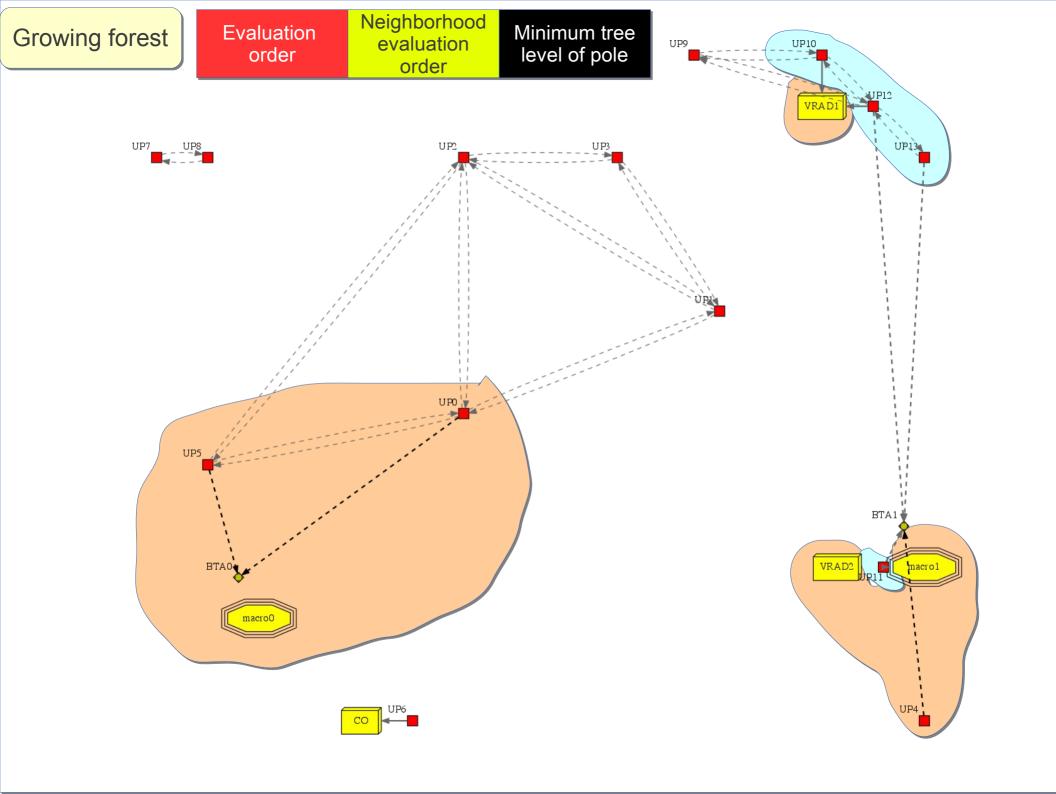


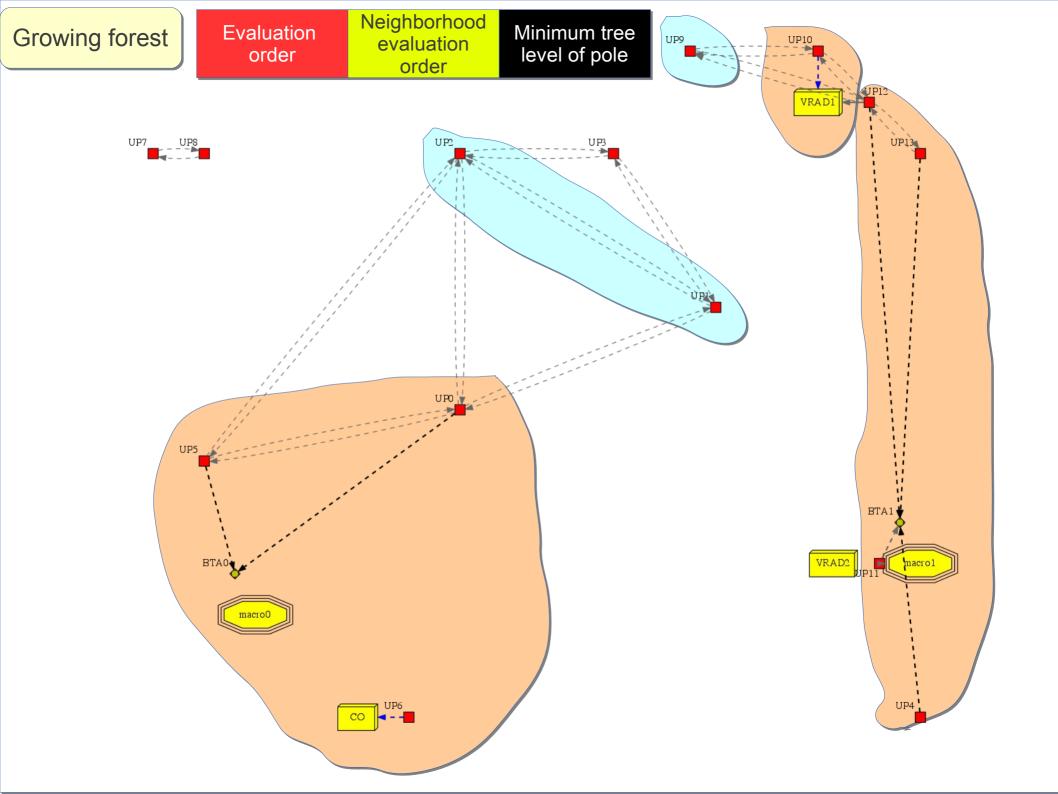


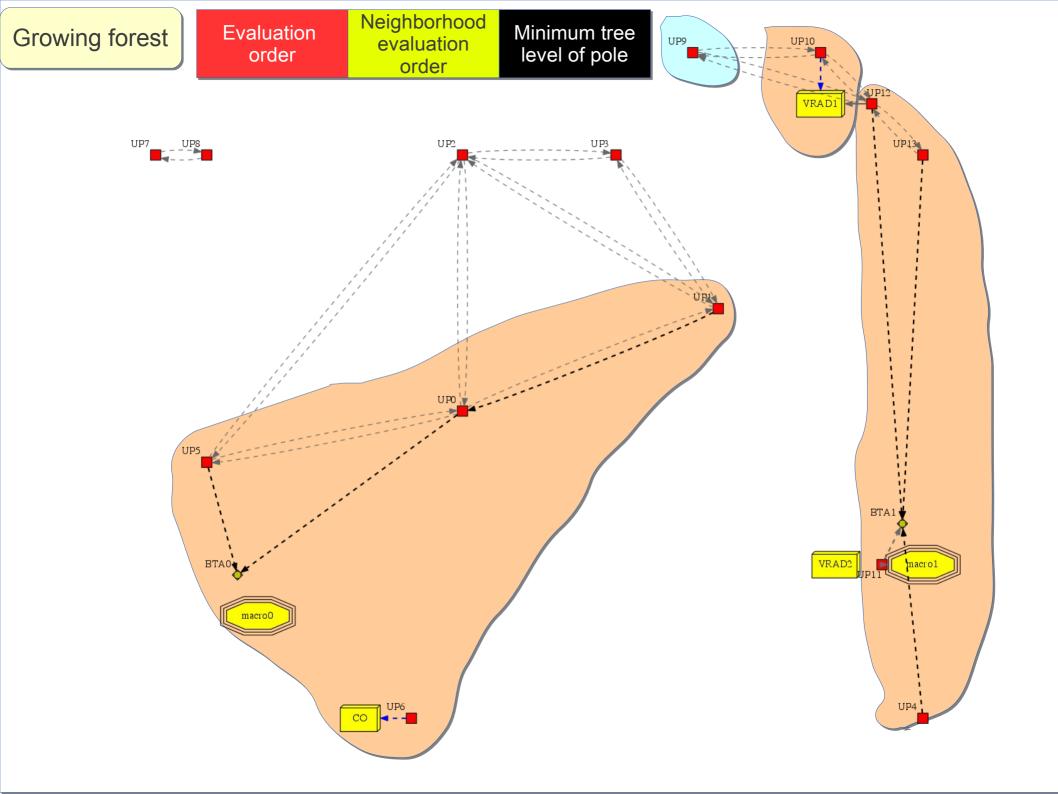


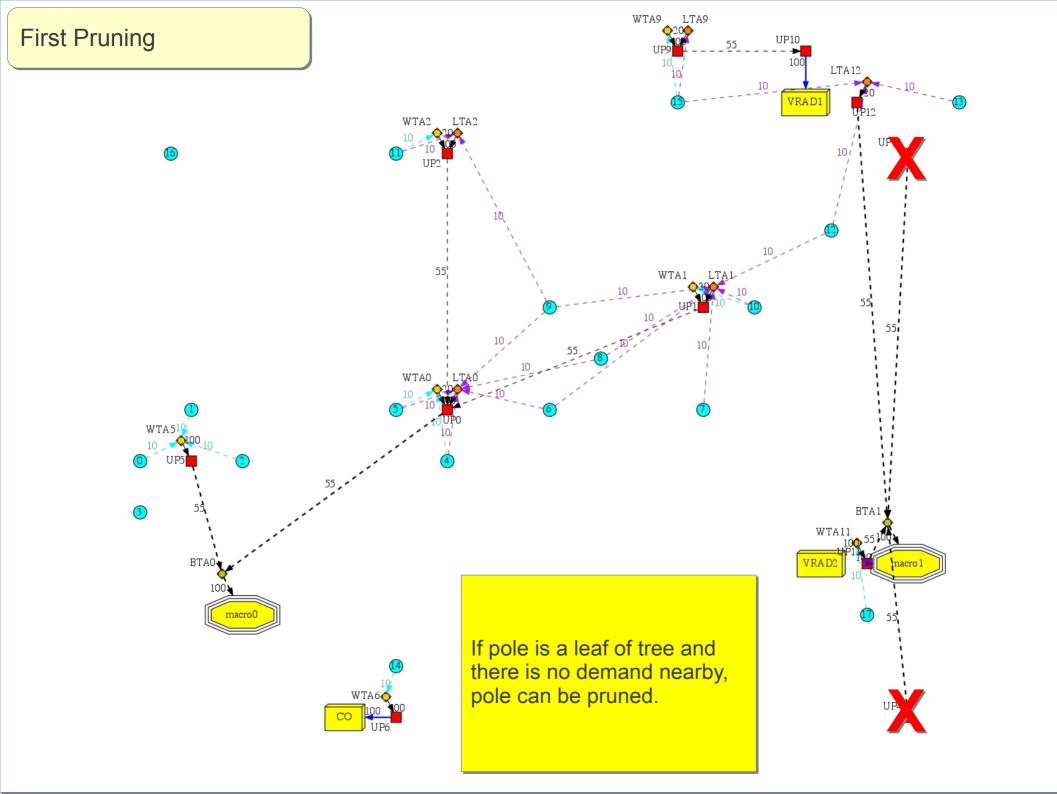


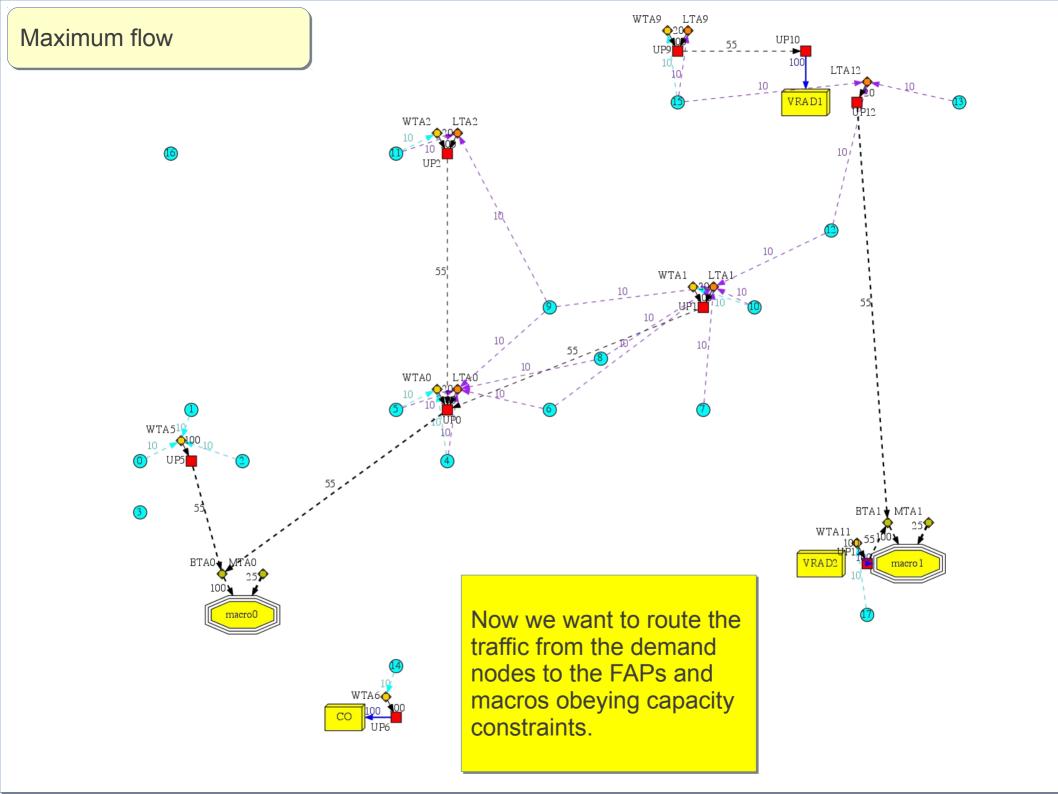


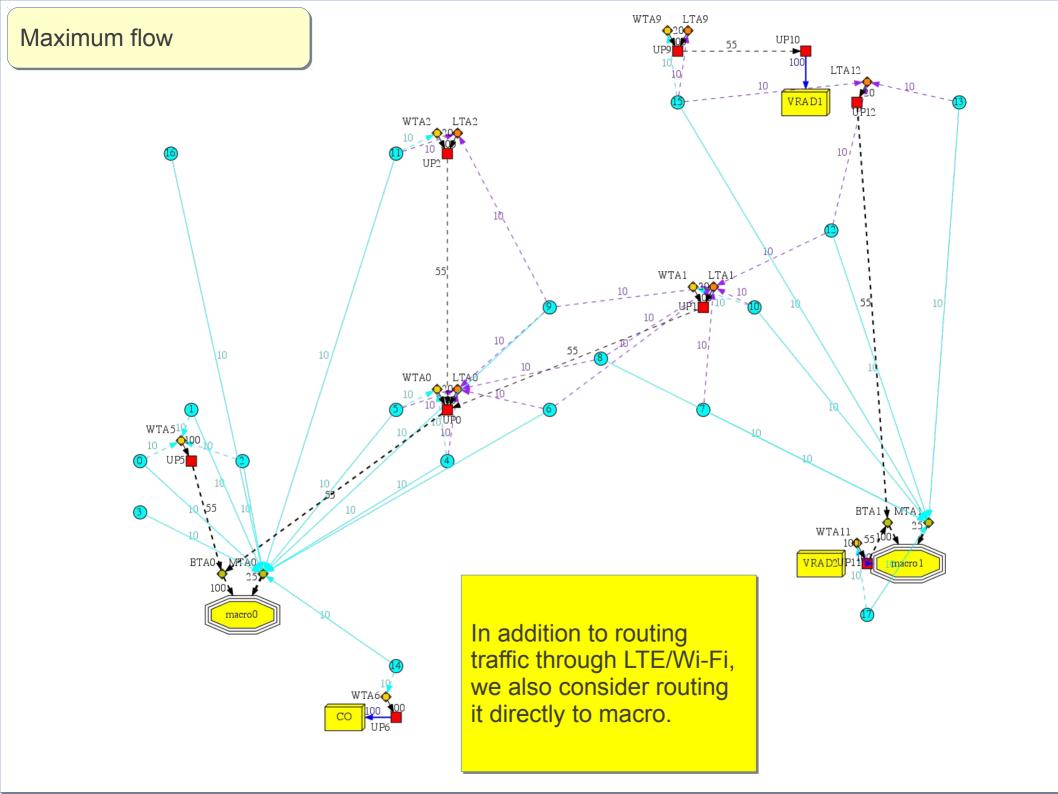


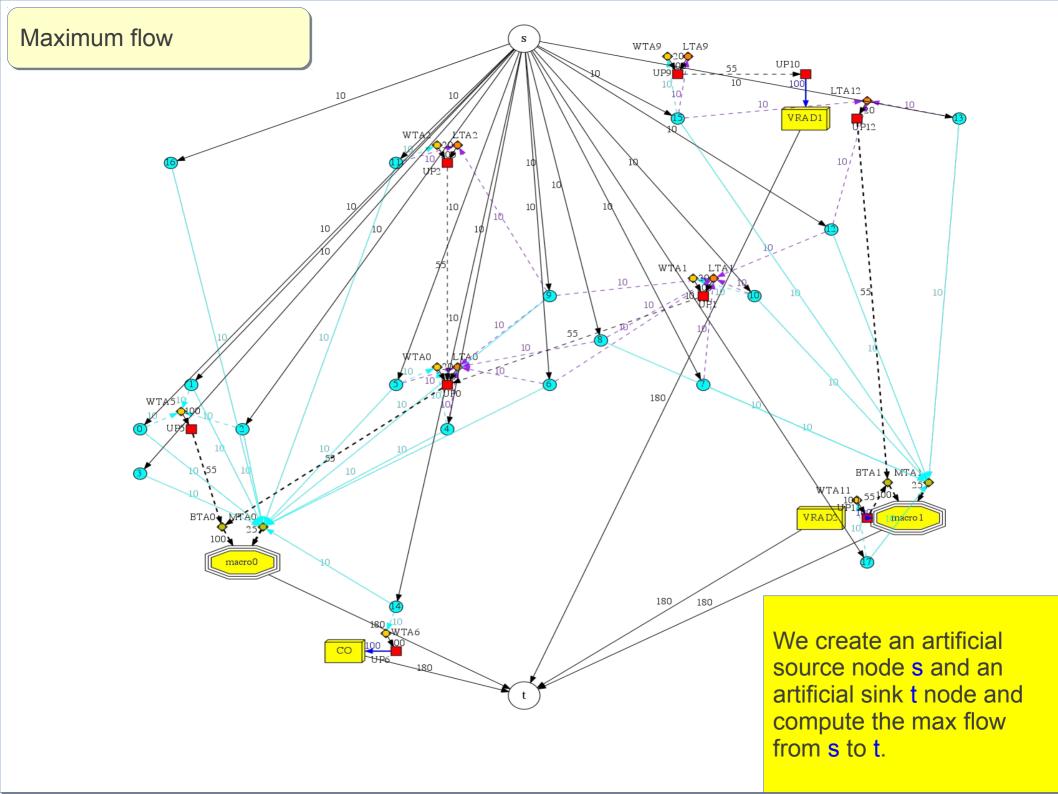


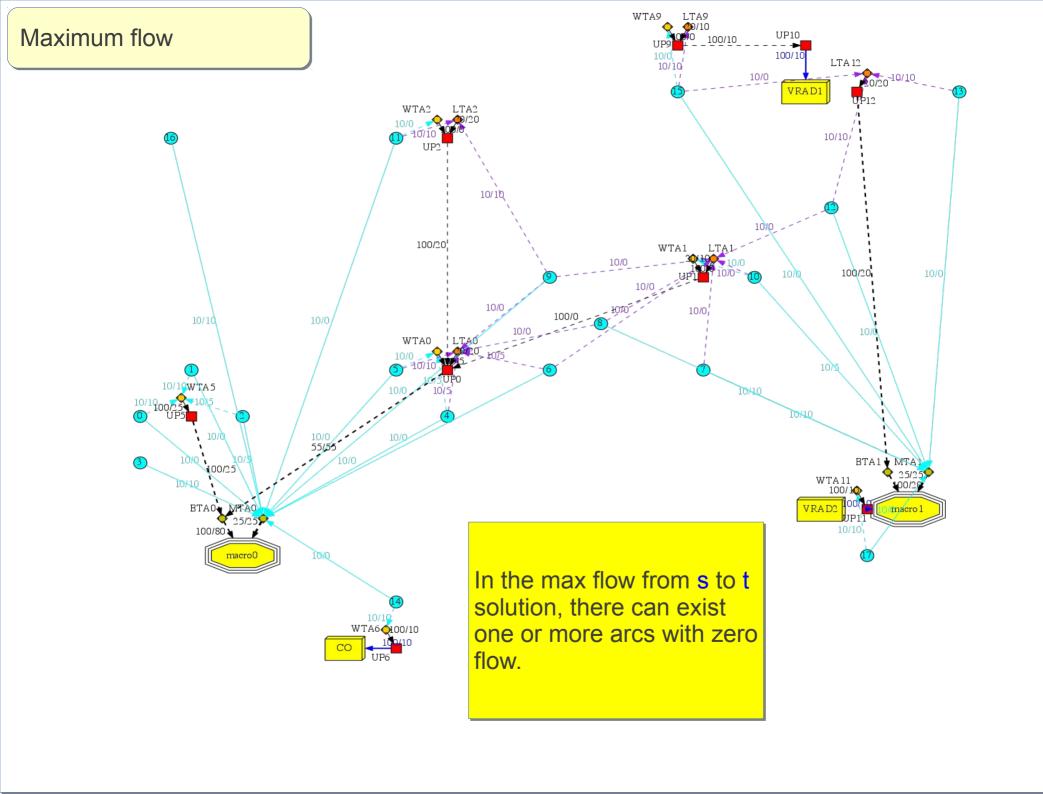


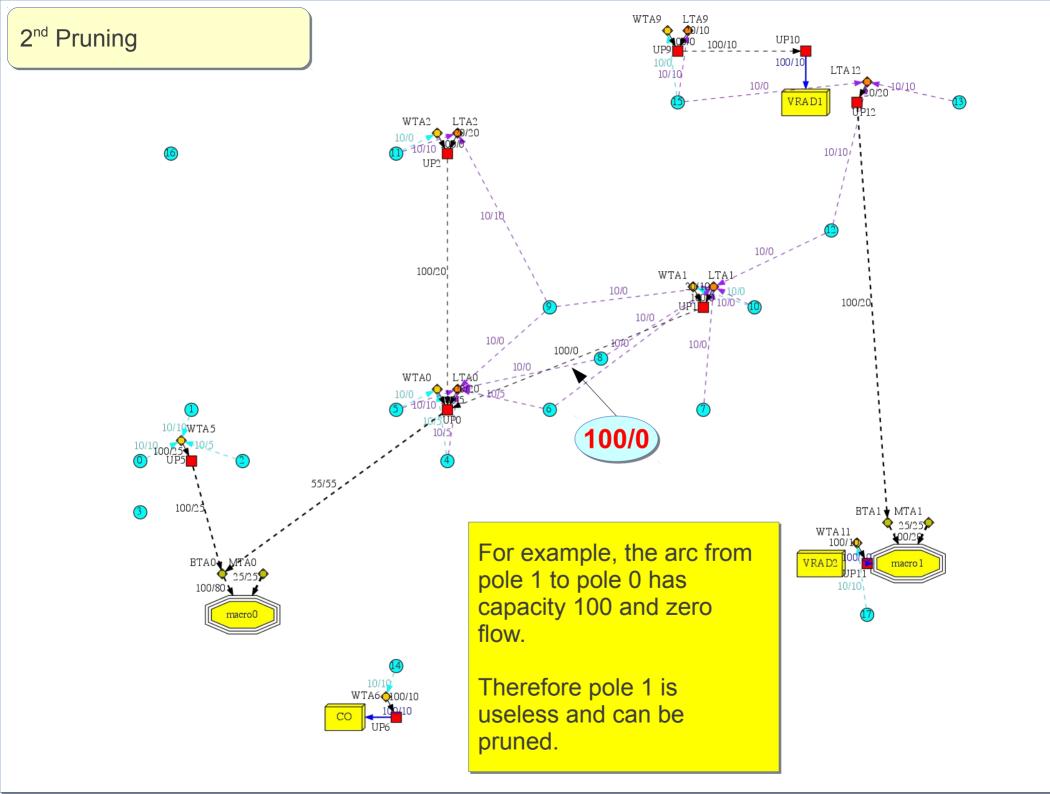


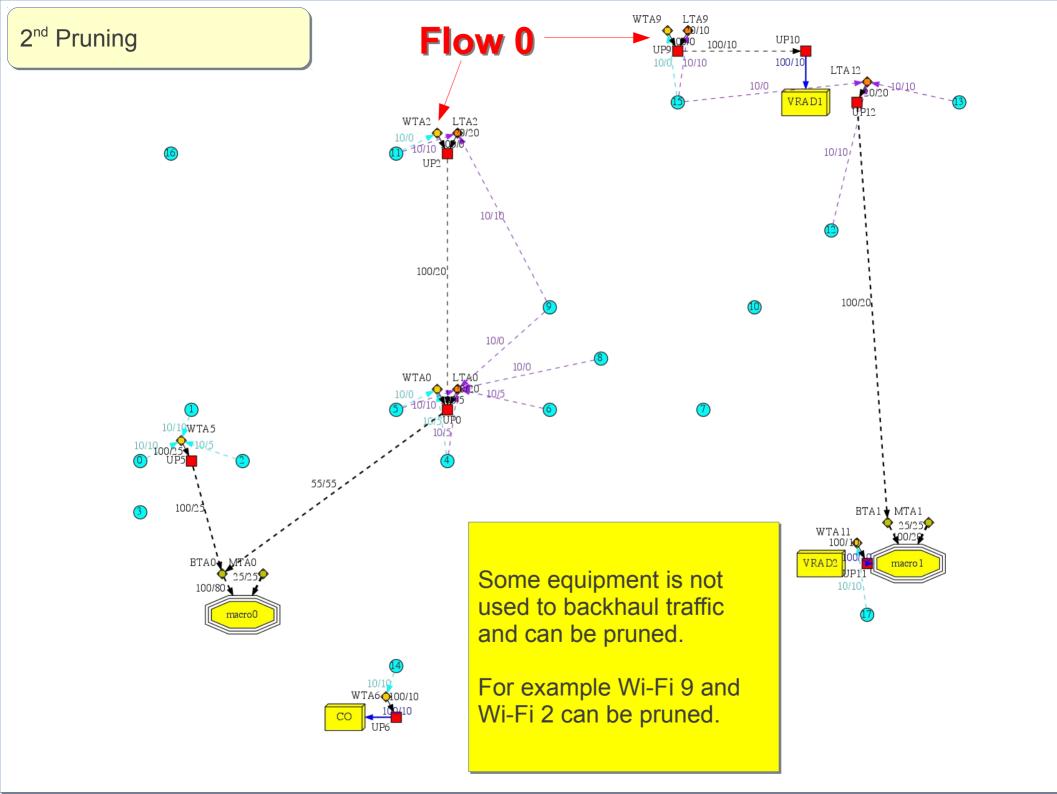




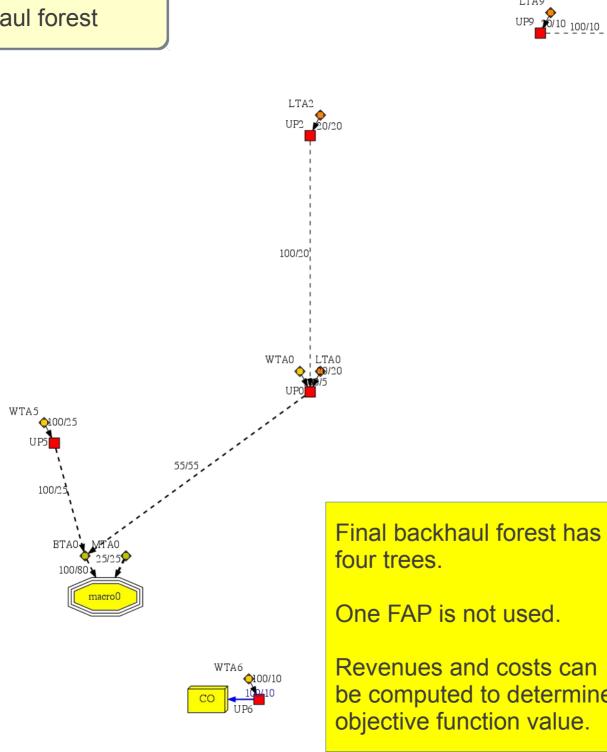


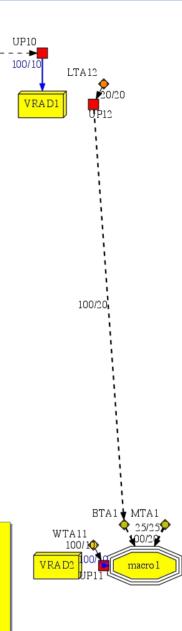






#### Final backhaul forest



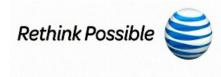


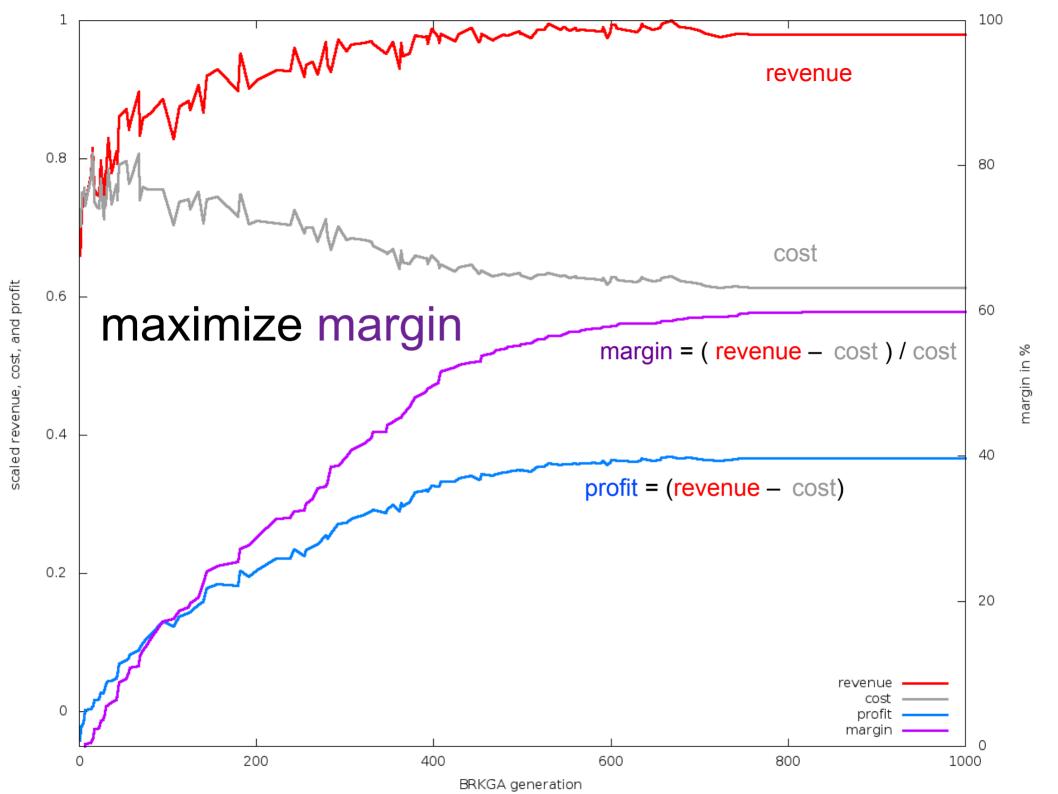
Revenues and costs can be computed to determine objective function value.

### An example optimization

We ran the BRKGA on a "real" instance from a large Tier 1 Internet Service Provider

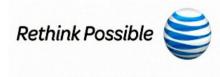
- Number of utility poles: 4395
- Number of demand points: 3236
- Number of fibered access points (FAP) excluding macrocells: 482
- Number of macrocells: 31
- Objective function: maximize margin = (revenue cost)/cost





### Concluding remarks

- We described the prize collecting directed k-hop Steiner forest (PCk-HSF) problem
- We modeled a wireless backhaul network planning problem as a PCk-HSF problem with additional constraints
- We described a biased random-key genetic algorithm (BRKGA)
  for the wireless backhaul network planning problem focusing on
  the decoder
- We applied the BRKGA to a "real" instance of the wireless backhaul network planning problem



# The End

These slides are available at http://mauricioresende.com

